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EXTERIOR: Using A Dual-VM Based External Shell for Guest-OS Introspection, Configuration, and Recovery

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March 17th, 2013

Our Approac

Evaluations

Conclusion





















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Virtualization



Virtualization Layer

Hardware Layer

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Virtualization



Virtualization Layer

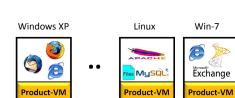
Hardware Layer

Virtualization (i.e., hypervisor) [Popek and Goldberg, 1974] has pushed our computing paradigm from multi-tasking to multi-OS.

Virtualization

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Virtualization Layer

Hardware Layer

Virtualization (i.e., hypervisor) [Popek and Goldberg, 1974] has pushed our computing paradigm from multi-tasking to multi-OS.

Consolidation, Migration, Isolation ...

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Execution Mode

Linux kill ps Product-VM

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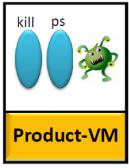
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Execution Mode

Linux



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Execution Mode

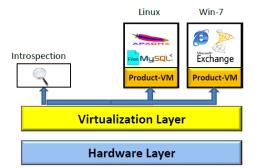
Linux

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Virtual Machine Introspection (VMI) [Garfinkel et al, NDSS'03]



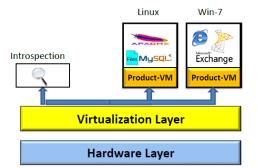
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Virtual Machine Introspection (VMI) [Garfinkel et al, NDSS'03]



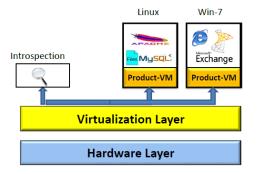
Using a trusted, dedicated virtualization layer program to monitor the running VMs

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Virtual Machine Introspection (VMI) [Garfinkel et al, NDSS'03]



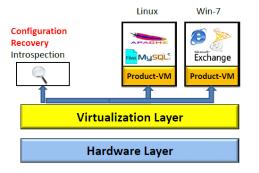
Using a trusted, dedicated virtualization layer program to monitor the running VMs

- Intrusion Detection
- Malware Analysis
- Memory Forensics

Evaluations

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Virtual Machine Introspection (VMI)



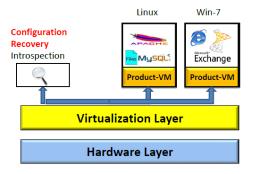
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Virtual Machine Introspection (VMI)



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EXTERIOR

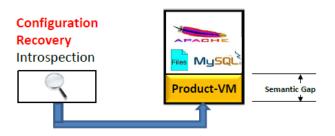
<u>Ex</u>ecute trusted utilities in SVM for <u>timely</u> Guest-OS introspection, (re)configu<u>ration</u> and <u>recovery</u>.

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The Semantic Gap in VMI ([Chen and Noble HotOS'01])



Linux

• View exposed by Virtual Machine Monitor is at low-level

- There is no abstraction and no APIs
- Need to reconstruct the guest-OS abstraction









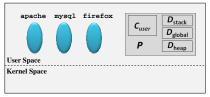


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Using a Dual-VM Architecture



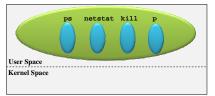
Guest VM (GVM)



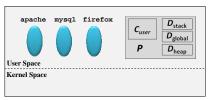
Our Approach ●00000 Evaluations

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Using a Dual-VM Architecture



Secure VM (SVM)



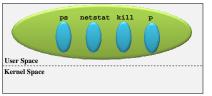
Guest VM (GVM)

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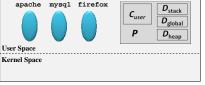
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Using a Dual-VM Architecture





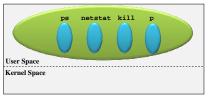


Guest VM (GVM)

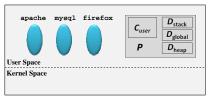
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- Virtual Machine Introspection
- Virtual Machine Configuration
- Intrusion Detection, Prevention (Recovery)

Overview	Our Approach	Evaluations	Conclusion
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Advantages			









- Isolation (SVM and GVM are isolated)
- Trustworthiness (trust code is running in secure VM)
- Automation (no need to develop introspection utilities)
- Security (enabling malware analysis, forensics...)
- **Transparency** (programmers write native program in SVM)

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Observ	ration		
1 e	execve("/sbin/sysctl",["sysctl",	"-w","kernel=1"],) = 0
2 k	ork(0)	$= 0 \times 604000$	
3 a	access("/etc/ld.so.nohwcap",F_OB	(x) = -1 ENOENT	
4 n	<pre>map(NULL, 8192, PROT_READ , -</pre>	-1,0) = 0x7f07b1749000	נ
5 a	access("/etc/ld.so.preload",R_OB	(x) = -1 ENOENT	
6 0	open("/etc/ld.so.cache", O_RDONI	LY) = 3	
_ · · ·			
	<pre>open("/proc/sys/kernel/randomize stat(3, {st_mode=S_IFREG 0644,</pre>	• • • • •) = 3

 49 mmap(NULL, 4096, PROT_READ|..., -1, 0) = 0x7f07b1748000

 50 write(3, "1\n", 2) = 2

 51 close(3) = 0

 ...

 57 exit_group(0) = ?

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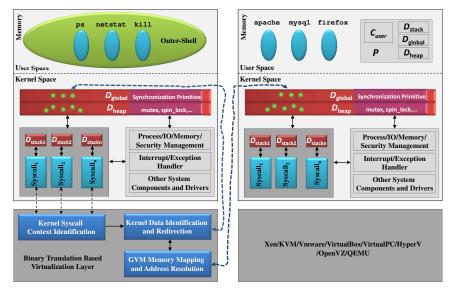
Syscall trace of running $\tt sysctl -w$ to turn on the address space randomization in Linux kernel 2.6.32

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Architecture Overview of EXTERIOR



Secure VM (SVM)

Guest VM (GVM)

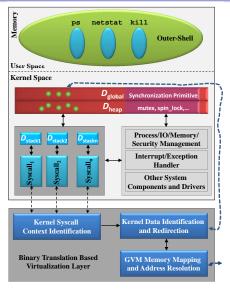
Overview

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The algorithms

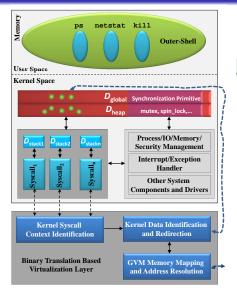


Secure VM (SVM)

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The algorithms



The Algorithm

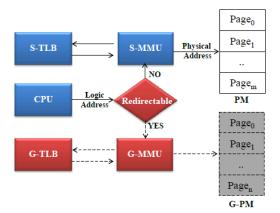
4.	
	DynamicBinaryInstrumentation(i):
2:	if SysCallExecContext(s):
3:	<pre>if SysCallRedirectable(s):</pre>
4:	RedirectableDataTracking(i);
5:	for $lpha$ in MemoryAddress(i):
6:	if DataRead(α):
7:	$P\!A(lpha) \leftarrow \mathtt{V2P}(lpha)$
8:	$Load(PA(\alpha))$
9:	else:
10:	if Configuration:
11:	$\mathtt{Store}(PA(\alpha))$
12:	else: //Introspection
13:	

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Mapping the GVM Memory Address



Outline

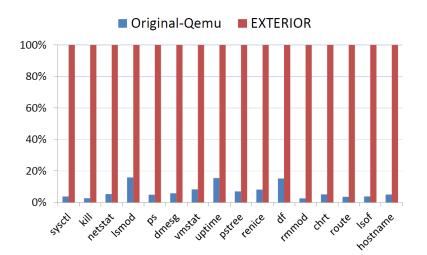
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Overview

Effectiveness

Cotogory	Utility	Effective?	
Category		Syntactics	Semantics
	ps (1)	X	\checkmark
	pstree (1)	X	\checkmark
	Ismod (8)	\checkmark	\checkmark
	dmesg (1)	\checkmark	\checkmark
Introspection	vmstat (8)	×	\checkmark
	netstat (8)	\checkmark	\checkmark
	lsof (8)	×	\checkmark
	uptime (1)	×	\checkmark
	df (1)	×	\checkmark
	sysctl (8)	\checkmark	\checkmark
Configuration	route (8)	\checkmark	\checkmark
	hostname (1)	\checkmark	\checkmark
	chrt (1)	\checkmark	\checkmark
	renice (1)	\checkmark	\checkmark
Bacovory	kill (1)	\checkmark	√
Recovery	rmmod (8)	\checkmark	\checkmark

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Overview ooooooo	Our Approach	Evaluations oo●oo	Conclusion
Recovery			

Rootkit	Targeted Function Pointer	Succeed?
adore-2.6	kernel global, heap object	×
hookswrite	IDT table	\checkmark
int3backdoor	IDT table	\checkmark
kbdv3	syscall table	\checkmark
kbeast-v1	syscall table, tcp4_seq_show	\checkmark
mood-nt-2.3	syscall table	\checkmark
override	syscall table	\checkmark
phalanx-b6	syscall table, tcp4_seq_show	\checkmark
rkit-1.01	syscall table	\checkmark
rial	syscall table	\checkmark
suckit-2	IDT table	\checkmark
synapsys-0.4	syscall table	\checkmark

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OS-Agnostic Testing

Linux Distribution	Kernel Version	Release Date	Transparent?
Debian 4.0	2.6.26	2007-04-06	\checkmark
Debian 5.0	2.6.28	2009-02-12	\checkmark
Debian 6.0	2.6.32	2010-01-22	\checkmark
Fedora-8	2.6.23	2007-11-08	 ✓
Fedora-10	2.6.27	2008-11-25	\checkmark
Fedora-12	2.6.31	2009-11-17	\checkmark
Fedora-14	2.6.35	2010-11-02	\checkmark
Fedora-16	3.1.0	2011-11-08	\checkmark
OpenSUSE-10.3	2.6.22	2007-10-04	✓
OpenSUSE-11.0	2.6.25	2008-06-19	\checkmark
OpenSUSE-11.1	2.6.27	2008-12-18	\checkmark
OpenSUSE-11.2	2.6.31	2009-11-12	\checkmark
OpenSUSE-11.3	2.6.34	2010-07-15	\checkmark
OpenSUSE-12.1	3.1.0	2011-11-16	√
Ubuntu-8.04	2.6.24	2008-04-24	√
Ubuntu-8.10	2.6.27	2008-10-30	\checkmark
Ubuntu-9.04	2.6.28	2009-04-23	\checkmark
Ubuntu-9.10	2.6.31	2009-10-29	√
Ubuntu-10.04	2.6.32	2010-04-29	✓
Ubuntu-10.10	2.6.35	2010-10-10	√
Ubuntu-11.04	2.6.38	2011-04-28	√
Ubuntu-11.10	3.0.4	2011-10-13	✓

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Limitations and Future Work

Limitations

- Can handle kernel ASLR
- Need an identical trusted kernel
- Need to stop the guest VM

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Evaluations

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Limitations and Future Work

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- Can handle kernel ASLR
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Future Work

- Derandomize the kernel address space
- Port to Windows OS

Outline

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Overview	Our Approach	Evaluations	Conclusion
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Conclusion			

• EXTERIOR is a novel dual-VM based external <u>shell</u> for trusted, native, out-of-VM program execution.

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Overview	Our Approach	Evaluations	Conclusion
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Conclusion			

- EXTERIOR is a novel dual-VM based external <u>shell</u> for trusted, native, out-of-VM program execution.
- It can be used for (automatic) introspection, (re)configuration of the guest-OS state (in the cloud), and can perform a timely response such as recovery from a kernel malware intrusion.

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Overview 0000000	Our Approach	Evaluations	Conclusion ●○
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- EXTERIOR has demonstrated a new program execution model on top of virtualization.

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Conclusion			

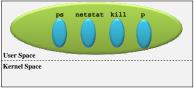
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- EXTERIOR has demonstrated a new program execution model on top of virtualization.
- (We believe) It will open new opportunities for system administration and security.

Overview

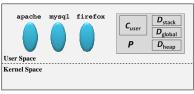
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Thank you !

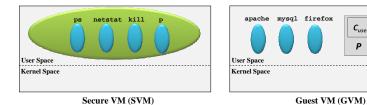


Secure VM (SVM)



Guest VM (GVM)

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Thank you !			



Contact us via. {yangchun.fu,zhiqiang.lin}@utdallas.edu for any questions

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Dglobal

Dheap

C_{user}

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